VISUAL OBJECT RECOGNITION

STATE-OF-THE-ART TECHNIQUES AND PERFORMANCE EVALUATION

LECTURE 6:TREE SEARCH AND HASHING

High dimensional spaces

*kD-Trees and Best Bin First (BBF)

Ball Trees

%K-means tree

Geometric Hashing

MOTIVATION

- * Finding the best match to a query descriptor q in a database with N prototypes p1...pN costs O(N).
- For a database with thousands or millions of descriptors this is expensive.
- A tree can find several good matches (near neighbours) in O(log N) time.
- A hash table can find *a good match* in O(1) time.

HIGH DIMENSIONAL SPACES

Distances in high dimensional spaces are higher on average!



HIGH DIMENSIONAL SPACES

Volume shrinks relative to area.
Example: unit "ball"

dimension	volume	area	volume/area
1	2r	2	r
2	πr^2	$2\pi r$	r/2
3	$4\pi r^{3}/3$	$4\pi r^2$	r/3
4	$\pi^{2}r^{4}/2$	$2\pi^2 r^3$	r/4

This means that a decision region in \mathbb{R}^D has increasingly more edges as D increases.

HIGH DIMENSIONAL SPACES

Box decision regions have 2 edges in R¹, 4 in R², 6 in R³, 8 in R⁴,... 2D in R^D



...

KD-TREES

A binary tree for search in \mathbb{R}^k



KD-TREES

Sinary search only works in 1D, in higher dimensions the kD-tree gives a *near neighbour*.

Tree construction algorithm:

 Select dimension k_n with largest variance
 Split dataset in two along selected dimension at median value, m_n.
 Repeat for each of the subsets.

KD-TREES

- Search for one neighbour is just one pass down the tree, and thus computation time is proportional to tree depth, d
- \ll Tree depth $d = \lceil \log_2 N \rceil$

* To find more neighbours, the original algorithm suggested a depth-first search with branch pruning.

If $e_{curr} < q[k_n]-m_n$ then skip branch.

BEST-BIN-FIRST

Depth-first search works poorly in high dimensional spaces, and thus Beis&Lowe suggest a best-first search instead.

% Algorithm:

At each node, store the distance
 e_n=q[k_n]-m_n in a *priority queue*. Always insert
 lowest value first.

2. Go down alternate branch of the first node in the queue if $e_n < e_{curr}$

BEST-BIN-FIRST





BALL TREES

Somohundro 89, Metric Tree Uhlmann 91

Second Secon

* Each node in tree has a centre p, and a radius r

p is average of all leaves r is maximum distance from p to a leaf

BALL TREES

- ** An optimal ball tree is constructed bottom up. Very expensive. E.g. using *agglomerative clustering*:
 - Set each sample to be one cluster
 Merge the two most similar clusters
 Repeat step 2 until no clusters are left.

* Agglomerative clustering generates a *dendrogram*, or similarity tree. This can be pruned using varius heuristics to form ball tree.

BALL TREES

Example of a search:



Leibe&Mikolajczyk&Schiele, BMVC'06

At each node, the distances to circle centres are computed, and compared to the radius.



Advantage: Good if range search is needed.
I.e. find all neighbours with d<d_{max}

Disadvantages:

Tree construction algorithm does not scale to very large datasets

* A ball in \mathbb{R}^D is not such a useful region shape if sample density varies in the feature space.

David Nistér and Henrik Stewénius, Scalable Recognition with a Vocabulary Tree, CVPR06

Hierarchical modification of the visual words idea from LE5



- Building the tree:
 - Run K-means with e.g. K=10 on whole dataset.
 - 2. Partition dataset into K subsets using Voronoi regions
 - 3. Apply algorithm recursively on subsets.

The tree gets branching factor K.

Wing the tree:

1. Compare query vector to prototypes at current level.

2. Go down best branch



Subset to compute a TF-IDF bag-of-words vector quickly.

Much faster than non-hierarchical visual words algorithm.

As in the kD-tree, the terminal leaf node is a near neighbour.

** Not so successful as a near neighbour search, as prototypes are far from region edge in \mathbb{R}^D .

HASH TABLES

An efficient way to perform lookup.



Seach key is converted to an index using a bashing function: index=H(key)

HASH TABLES

Lookup is O(1) instead of e.g. O(N) in a list, O(log N) in a sorted list/tree etc.

Collisions can happen. i.e. different keys get the same index. Solved e.g. using *chaining* (linked lists), or *linear probing* (insertion at next free slot).

Linear probing typically wants a <80% filled table.</p>

Hashing has poor cache locality.

GEOMETRIC HASHING

Introduced in Lamdan&Wolfson ICCV'88



GEOMETRIC HASHING

Modern example: Used for matching frames without descriptors by Chum & Matas, Geometric Hashing with Local Affine Frames, CVPR'06

We pairs of affine frames. Express frame 2 in frame 1. 25 bins for angle 16 for d₁, 6 for d₂ & d₃

\$\$ 9*10⁶ unique values for key to hash.



GEOMETRIC HASHING

Design of H(key) is not discussed further.

* Hash tables suffer from the same basic problem as trees: Neighbouring bins might contain the closest match. More neighbours in high dimensional spaces.

To deal with the neighbour problem, Chum&Matas construct 6 different tables (for 6 different frame constructions) and run them in parallel.

PROJECTS

Course is 8hp: 5hp for lectures+articles+exam 3hp for project.

Project part is 2 weeks of: programming&research writing a small report.

PROJECT SUGGESTIONS

1.Comparison of kD-tree and K-means tree in terms of speed and accuracy.

2.Implementation of a bag-of-words recognition system (using existing feature detector code). Test how system parameters affect result.

3.Implementation and test of a new descriptor for a given detector.

PROJECT SUGGESTIONS

4.Implement and test a voting scheme for global geometric deformation (e.g. similarity, or affine transform) of feature locations.

5.Learn a matching metric, and compare it to least squares matching.

6.Compare Chi², EMD and least-squares on a problem of choice.

7. Your own suggestion.



Everyone should bring calendar next week, so we can decide on a date for the written exam.

DISCUSSION

Questions/comments on paper and lecture.